

Towards Generic Countermeasures Against Fault Injection Attacks

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- ▶ Allows to recover the secret primes p and q used in the secret keys of the CRT-RSA cryptosystem.
- ▶ Only requires a single fault injection and a gcd computation.
- Many countermeasures have been developed.

- ▶ Mostly resulting from engineering efforts.
- ▶ Development by trial-and-error leading to overkill protections.
- ▶ Many different countermeasures (NIH, patents), not all of them work.

- ▶ Formal studies of these countermeasures allowed to understand their working factor.
- We were able to fix the broken ones and to simplify many of them (e.g., original Vigilant's countermeasure: broken, 9 tests, 5 random numbers; our fixed and simplified version: working, 3 tests, 1 random number).
- ▶ More importantly, the working factor is actually not tied to the BellCoRe attack, nor to the CRT-RSA algorithm.
- It is possible to abstract it and get a recipe for cost-effectively verifying the integrity of any arithmetic computation.

- ▶ Idea: verify the integrity of the computation by introducing redundancy.
- ▶ Simply repeating the computation and comparing results is bad:
 - (a) it is too expensive, and
 - (b) nothing stops the attacker from injecting the same fault twice.
- ▶ Thus, existing countermeasures optimize this idea in different ways.

- ▶ The *entanglement* protection scheme solves both issues, by:
 - ▶ lifting the computation to an over-structure (a direct product) allowing
 - (a) to project the result back onto the original structure, and
 - (b) to project a checksum onto a smaller structure (e.g., `int32`-sized);
 - ▶ performing in parallel the same computation in the smaller structure;
 - ▶ both the checksum and the smaller result should be equal.
- ▶ The redundant part of the computation is almost free (arithmetic with 32-bit vs. 2,048-bit numbers).
- ▶ It is very hard to precisely fault the small computation to produce a consistent value modification.
- ▶ Limitation: possible collisions in the small structure.
Mitigated by the possibility to use several different small structures.

- ▶ At IMDEA Software Institute (Madrid, Spain), I developed a compiler called enredo, while supervised by Gilles Barthe, François Dupressoir and Pierre-Yves Strub.
- ▶ Automated insertion of the *entanglement* countermeasure into arbitrary code.
- ▶ ~~Short demo.~~

```
Field  $\mathbb{Z}_p$  ;
```

```
 $\mathbb{Z}_p$  a, b, c, d ;
```

```
a := b + (c * d) ;  
return a ;
```



```
Field Zp, Zr0 ;
Ring Zpr0 = Zp * Zr0 ;

Zp _ret_a ;
Zr0 _ar0, _br0, _cr0, _dr0 ;
Zpr0 a, b, c, d ;

_dr0 := d [Zr0] ;
_cr0 := c [Zr0] ;
_br0 := b [Zr0] ;
_ar0 := _br0 + (_cr0 * _dr0) ;
a := b + (c * d) ;
if a [Zr0] = _ar0
then
  _ret_a := a [Zp] ;
  return _ret_a ;
else
  return ERROR ;
end
```

```

Field Zp, Zr0, Zr1 ;
Ring Zpr0r1 = Zp * Zr0 * Zr1 ;

Zp _ret_a ;
Zr0 _ar0, _br0, _cr0, _dr0 ;
Zr1 _ar1, _br1, _cr1, _dr1 ;
Zpr0r1 a, b, c, d ;

_dr1 := d [Zr1] ;
_cr1 := c [Zr1] ;
_br1 := b [Zr1] ;
_dr0 := d [Zr0] ;
_cr0 := c [Zr0] ;
_br0 := b [Zr0] ;
_ar0 := _br0 + (_cr0 * _dr0) ;
_ar1 := _br1 + (_cr1 * _dr1) ;
a := b + (c * d) ;
if a [Zr1] = _ar1 /\ a [Zr0] = _ar0
then
  _ret_a := a [Zp] ;
  return _ret_a ;
else
  return ERROR ;
end

```

- ▶ We already have:
 - ▶ an executable code output (Python),
 - ▶ a correctness proof of the code transformation.
- ▶ Benchmark of the cost of the countermeasure.
- ▶ Security proof.
- ▶ Protected implementations of currently unprotected algorithms.
- ▶ Practical lab tests.
- ▶ This work will appear in a chapter of *Handbook of Pairing Based Cryptography* (Editors: Nadia El Mrabet and Marc Joye).

The BellCoRe Attack
State-of-the-Art Countermeasures
Formal Study of Countermeasures
Integrity Verification
Entanglement
enredo
Perspectives

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