Capacity: an Abstract Model of Control over Personal Data

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Paper (with Daniel Le Métayer): hal-01638190

- Control over personnal data
- Modeling control
- Characterizing control
- Evaluating concrete systems

- > The notion of *privacy by control* is predominant in the privacy literature.
- However, it lacks a formal definition.
- > This makes it hard to check for compliance, to compare design options, etc.
- → We want a formal framework to specify the notion of *control over personal data*.

- Formally capturing the notion of control is notoriously difficult.
- Control is about a potential rather than one particular realization.
- Existing control literature (e.g., access control and usage control) does not really encapsulates the intuition underlying the notion of control over personal data.

- In their 2015 paper*, Lazaro and Le Métayer identified three dimensions of control over personal data.
- > These three dimensions corresponds to the capacities for an individual:
 - to perform actions on their personal data,
 - to prevent others from performing actions on their personal data, and
 - to be informed of actions performed by others on their personal data.
- → Based on this work, we built *Capacity*.

* http://script-ed.org/?p=1927

- Capacity's goal is to model control over personal data in a very general way.
- Thus, guiding principles of its design are abstraction and minimality.
- Basically, agents can perform operations on resources in given contexts.
- Control is modeled by requirements expressing constraints on those operations.
- → Running example for this: rudimentary photo sharing service.

- > This talk uses a simple photo sharing service, named Album, as an example.
- Album is a centralized service where:
 - users can upload, delete, and access photos in their album ;
 - users can connect to each other to become friends ;
 - users can see their friends photos ;
 - users can tag theirs and their friends photos with their name or the names of friends ;
 - users are notified when they are tagged in a photo by someone else.

- There are four types of atomic objects in Capacity:
 - Agents:
 - agents model users and services,
 - the set of agents is \mathcal{A} ,
 - examples: Album (the service) and its users (Daniel, Pablo, ...);

Resources:

- resources model data, and typically personal data,
- the set of resources is ${\cal R}$,
- examples: usernames (Pablo), users' album (*album*_{Pablo}), and photos (💹).

Operations:

- operations model what can be performed on resources,
- the set of operations is \mathcal{O} ,
- examples: connect, upload, tag, access, delete ;

Contexts:

- contexts model any external factors relevant to an operation,
- the set of contexts is \mathcal{C} ,
- examples: location, time, relationship between agents, purpose, exposure.

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> Actions model the application of an operation to a list of parameters in a context.

- Action $op_c(x_1, \ldots, x_n)$ is the application of operation op to x_1, \ldots, x_n in context c.
- Parameters x_i can be resources or agents.
- Examples:
 - connect_c(Daniel),
 - $upload_c(\mathbb{M}, album_{Pablo})$,
 - $tag_c(\mathbb{Z}, Daniel)$.
- The set of actions is Δ.

- ▶ We define three *relations* on atomic objects:
 - Pers(r, a) expresses that resource r is a personal data of agent a,
 - $In(r, \alpha)$ expresses that resource r is involved in action α ,
 - *Trust*(*a*, *b*) expresses that agent *a* trusts agent *b*.
- ► Examples:
 - *Pers*(ᢂ, Pablo),
 - In(■, tag_c(■, Pablo)),
 - *Trust*(Pablo, Daniel).

- A requirement R is a relation $Can^R \subseteq \mathcal{A} \times \Delta \times \mathcal{P}(\mathcal{A}) \times \mathcal{P}(\mathcal{A})$.
- Intuitively, $Can^{R}(a, \alpha, E, W)$ means that:
 - agent a can perform action α
 - only if this action is enabled by all agents in *E*
 - while all agents in Whave to to be informed of it.
- Examples:
 - $Can^{R}_{-}(Pablo, upload_{c}(\blacksquare, album_{Pablo}), \{Album\}, \{Album\}),$
 - $Can^{R}(\text{Daniel}, \text{upload}_{c}(\blacksquare, album_{Pablo}), \{\bot\}, \{\bot\}),$
 - $Can^{R}(Pablo, tag_{c}(\blacksquare, Daniel), \{Daniel, Album\}, \{Daniel, Album\}).$

> This single relation can express the three capacities of control of personal data:

- when x = a it expresses the capacity of x to *perform* action α ,
- when $x \in E$ it expresses the capacity of x to *prevent* action α ,
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- Requirements semantics is given by characterizing execution traces.
- Traces are characterized using four abstract properties:
 - $\theta \vdash \mathsf{Requests}(a, \alpha)$:
 - in trace θ , agent *a* attempts to perform action α ,
 - example: $\theta \vdash Requests(Pablo, tag_c(M, Daniel));$
 - $\theta \vdash \text{Enables}(a, b, \alpha)$:
 - in trace θ , agent a enables the performance of action α by agent b,
 - example: $\theta \vdash Enables(Album, Pablo, tag_c(M, Daniel)),$ $\theta \vdash Enables(Daniel, Pablo, tag_c(M, Daniel));$
 - $\theta \vdash \text{Does}(a, b, \alpha)$:
 - in trace heta, agent a performs action lpha on behalf of agent b,
 - example: $\theta \vdash Does(Album, Pablo, tag_c(M, Daniel));$
 - $\theta \vdash Notifies(a, b, c, \alpha)$:
 - in trace θ , agent a notifies to agent b the performance of action α on behalf of agent c,
 - example: heta dash Notifies(Album, Daniel, Pablo, $ag_c(\blacksquare,$ Daniel)).

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- A trace θ is consistent if:
 - $\theta \vdash \text{Does}(c, a, \alpha) \implies \theta \vdash \text{Requests}(a, \alpha),$
 - $\theta \vdash \mathsf{Notifies}(a, b, c, \alpha) \implies \exists d, \theta \vdash \mathsf{Does}(d, c, \alpha).$
- Intuitively, a trace is inconsistent if it includes:
 - an action performed on behalf of an agent that has not requested it, or
 - the notification of an action that has not been performed.

- A trace θ is *complete* wrt requirement R where $Can^{R}(a, \alpha, E, W)$ if:
 - $\theta \vdash \mathsf{Requests}(a, \alpha) \land \forall b \in E, \theta \vdash \mathsf{Enables}(b, a, \alpha) \implies \exists c \in \mathcal{A}, \theta \vdash \mathsf{Does}(c, a, \alpha).$
- > Intuitively, a trace is complete if an action is always performed when:
 - it has been requested, and
 - it has been enabled by all necessary agents.

- A trace θ is *compliant* with requirement R where $Can^{R}(a, \alpha, E, W)$ if:
 - $\forall d \in \mathcal{A}, \theta \vdash \mathsf{Does}(d, a, \alpha) \implies \forall b \in E, \theta \vdash \mathsf{Enables}(b, a, \alpha),$
 - $\forall d \in \mathcal{A}, \theta \vdash \mathsf{Does}(d, a, \alpha) \implies \forall b \in W, \exists c \in \mathcal{A}, \theta \vdash \mathsf{Notifies}(c, b, a, \alpha).$
- Intuitively, a trace is compliant if all Can^R constraint are met:
 - no action is performed unless it is enabled by all its enablers, and
 - all agents that have to be informed are notified.
- Compliance is noted $\theta \models R$.

▶ We introduce four independent *types of control*:

- action control,
- observability control,
- authorization control,
- notification control.
- Each type comes with three levels of control:
 - absolute control,
 - relative control,
 - lack of control.

- Action control describes an agent's control on actions that it initiates.
- With regard to a requirement R, an agent a has:
 - *absolute action control* over α if it does not depend on others to perform it:
 - $AA_R(a, \alpha) \iff Can^R(a, \alpha, \varnothing, W)$;
 - relative action control over α if it depends only trusted agents:
 - $RA_R(a, \alpha) \iff Can^R(a, \alpha, E, W) \land b \in E \implies Trust(a, b).$
- Examples:
 - $Trust(Pablo, Album) \implies RA_R(Pablo, upload_c(\blacksquare, album_{Pablo})),$
 - $Trust(Pablo, Album) \implies RA_R(Pablo, delete_c(\blacksquare, album_{Pablo})).$

- Observability control describes an agent's capacity to perform actions that are not observable by others.
- With regard to a requirement R, an agent a has:
 - absolute observability control over α if it can perform α discreetly:
 - $AO_R(a, \alpha) \iff Can^R(a, \alpha, E, \emptyset)$;
 - relative observability control over α if only trusted agents can know about it:
 - $RO_R(a, \alpha) \iff Can^R(a, \alpha, E, W) \land b \in W \implies Trust(a, b).$
- Examples:
 - $Trust(Pablo, Album) \implies RO_R(Pablo, upload_c(\blacksquare, album_{Pablo})),$
 - $Trust(Pablo, Album) \land Trust(Pablo, Daniel) \implies RO_R(Pablo, tag_c(\blacksquare, Daniel)).$

- Authorization control describes an agent's control on actions initiated by others.
- With regard to a requirement R, an agent a has:
 - absolute authorization control over α if it is the only agent that can prevent it:
 - $AH_R(a, \alpha) \iff Can^R(b, \alpha, \{a\}, W)$;
 - relative authorization control over α if it is not the only agent having this capacity:
 - $RH_R(a,\alpha) \iff Can^R(b,\alpha,E,W) \implies a \in E.$
- Examples:
 - $AH_R(\text{Album}, \text{upload}_c(\mathbb{M}, album_{\text{Pablo}}))$,
 - $RH_R(\text{Daniel}, \text{tag}_c(\mathbb{M}, \text{Daniel})).$

- Notification control describes an agent's capacity to be informed about actions performed by others.
- With regard to a requirement R, an agent a has:
 - absolute notification control over α if it is the only agent that has the capacity to be informed of it:
 - $AN_R(a, \alpha) \iff Can^R(b, \alpha, E, \{a\})$;
 - relative notification control over α if it is not the only agent having this capacity:
 - $RN_R(a, \alpha) \iff Can^R(b, \alpha, E, W) \implies a \in W.$
- Examples:
 - $AN_R(\text{Album}, \text{upload}_c(\textbf{M}, album_{Pablo}))$,
 - $RN_R(Daniel, tag_c(\mathbb{M}, Daniel)).$

- These types of control can be extended to resources and agents:
 - for resources, by generalizing to all actions that involves the resource:
 - e.g., $AA_R(a, r) \iff \forall \alpha \in \Delta, In(r, \alpha) \implies AA_R(a, \alpha)$;
 - for agents, by generalizing to all the personal data of the agent:
 - e.g., $AA_R(a) \iff \forall r \in \mathcal{R}, Pers(r, a) \implies AA_R(a, r).$
- Control lattice:
 - it is easy to check that absolute control implies relative control;
 - using the order defined by implication, we have a lattice made of 3⁴ forms of control for each action, data, and agent.

- > Concrete traces are sequences of concrete events which can be clearly identified:
 - HTTP requests and responses, SQL queries, file manipulations, etc.
- Modeling a concrete system in Capacity requires to:
 - identify the sets of agents, resources, actions, and contexts ;
 - define the conditions under which a concrete trace satisfies each abstract trace property.
- Given this model it is possible to:
 - compute the requirement that corresponds to the system,
 - verify if the system satisfies a specific requirement,
 - evaluate the types and levels of control of each agents.

- ▶ In Album, concrete traces are sequences of the following events:
 - U-registers(u): user u creates an account on Album ;
 - U-uploads-pic(u, p): user u uploads a photo to their album ;
 - U-requests-album(u_1, u_2): user u_1 requests u_2 's album;
 - U-submits-tag(u_1, p, u_2): user u_1 tags u_2 in photo p;
 - U-deletes-pic(u_1, p): user u_1 deletes photo p from their album ;
 - U-requests-con(u_1, u_2): user u_1 requests to connect with u_2 ;
 - U-accepts-con(u_1, u_2): user u_1 accepts to connect with u_2 ;
 - U-rejects-con(u_1, u_2): user u_1 rejects to connect with u_2 ;
 - U-disconnects(u_1, u_2): user u_1 disconnects from u_2 ;
 - A-creates-account(*u*): Album creates *u*'s account;
 - A-publishes-pic(p, u): Album publishes photo p in u's album ;
 - A-serves-album(u_1, u_2): Album sends u_1 the album of u_2 ;
 - A-connects(u_1, u_2): Album connects u_1 and u_2 ;
 - A-disconnects(u_1, u_2): Album disconnects u_1 and u_2 ;
 - A-tags-pic(u_1, p): Album tags u_1 in photo p;
 - A-notifies-req (u_1, u_2) : Album notifies u_2 of u_1 's request to connect;
 - A-notifies-con(u_1, u_2): Album notifies u_1 and u_2 that they are connected ;
 - A-notifies-tag(u_1, p, u_2): Album notifies u_1 that they have been tagged in photo p by u_2 .

- Let θ_n be the *n*th event in the concrete trace.
- We define our abstract properties as follows:
 - $\theta \vdash \text{Requests}(u, \text{upload}_n(p, album_u))$ $\iff \exists m < n, \theta_m = \text{U-uploads-pic}(u, p).$
 - $\theta \vdash Enables(\text{Album}, u, upload_n(p, album_u))$ $\iff \exists m < n, \theta_m = \text{U-registers}(u).$
 - $\theta \vdash Does(Album, u, upload_n(p, album_u))$
 - $\iff \theta_n = \text{A-publishes-pic}(p, u).$

- ▶ With these definitions we can prove that $\theta \models R$ such that:
 - $Can^{R}(u, upload_{n}(p, album_{u}), \{Album\}, \{Album\}).$
- Which in terms of control means that we have:
 - $RA_R(u, upload_n(p, album_u))$ if Trust(u, Album).
 - $RO_R(u, upload_n(p, album_u))$ if Trust(u, Album).
 - $AH_R(\text{Album}, \text{upload}_n(p, album_u)).$
 - $AN_R(\text{Album}, \text{upload}_n(p, album_u)).$

- Let θ_n be the *n*th event in the concrete trace.
- We define our abstract properties as follows:
 - $\theta \vdash Requests(u_1, tag_n(p, u_2))$

 $\iff \exists m < n, heta_m = \texttt{U-submits-tag}(u_1, p, u_2).$

• $\theta \vdash \text{Enables}(u_2, u_1, \text{tag}_n(p, u_2))$

 $\iff \exists m < n, heta_m = \texttt{U-accepts-con}(u_2, u_1) \lor heta_m = \texttt{U-accepts-con}(u_1, u_2)$

 $\land \nexists k, m < k < n, heta_k = U$ -disconnects(u_2, u_1) $\lor heta_m = U$ -disconnects(u_1, u_2).

- $\theta \vdash Enables(Album, u_1, tag_n(p, u_2))$ $\iff \theta_n = A \text{-tags-pic}(u_2, p).$
- $\theta \vdash \text{Does}(\text{Album}, u_1, \text{tag}_n(p, u_2))$ $\iff \theta_n = \text{A-tags-pic}(u_2, p).$
- $\theta \vdash Notifies(Album, u_2, u_1, tag_n(p, u_2))$
 - $\iff \theta_{n+1} = \text{A-notifies-tag}(u_2, p, u_1).$

- ▶ With these definitions we can prove that $\theta \models R$ such that:
 - $Can^{R}(u_{1}, tag_{n}(p, u_{2}), \{u_{2}, Album\}, \{u_{2}, Album\}).$
- Which in terms of control means that we have:
 - $RA_R(u_1, tag_n(p, u_2))$ if $Trust(u_1, Album) \wedge Trust(u_1, u_2)$.
 - $RO_R(u_1, tag_n(p, u_2))$ if $Trust(u_1, Album) \wedge Trust(u_1, u_2)$.
 - $RH_R(u, \operatorname{tag}_n(p, u)).$
 - $RN_R(u, \operatorname{tag}_n(p, u)).$
 - $RH_R(\text{Album}, ag_n(p, u)).$
 - $RN_R(\text{Album}, ag_n(p, u)).$

- Types and levels of control allow to formally compare different systems.
- Studying alternative implementations of a given specification can be useful for privacy by design.

- Capacity provides a formal framework to reason about privacy in terms of control.
- ▶ The goal of this work is to serve as foundation for new privacy research and tools.
- ► Future work (Tupac):
 - find a better way than contexts to formally capture the notion exposure ;
 - make a user-friendly interface to specify requirements ;
 - model control aspects of personal data related laws such as the GDPR ;
 - build tools to automate requirement verification.

Control over Personal Data

Control Three dimensions of control

Modeling Control with Capacity

Running example: Album Objects Actions Relations Requirements Abstract trace properties

Characterizing Control with Capacity

Action control Observability control Authorization control Notification control Extensions

Evaluating Concrete Systems with Capacity

An example with Album Album: uploading a photo Album: tagging a friend Implementations comparison

Conclusions and Perspectives